

5. Hamiltonian simulation (LCU): pt.1

Linear combination of Unitaries

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Recap

e^{-iHt}

- QPE is very useful for
 - ❖ energy estimation
 - ❖ ground state preparation
- QPE cost $\approx \underbrace{1/\epsilon}$ times Hamiltonian simulation cost for a unit time.
- Better Hamiltonian simulation \rightarrow Better time-evolution simulation, energy estimation, ground state preparation, ...
- Trotter-Suzuki: An efficient algorithm for Hamiltonian simulation on a quantum computer

Post-Trotter methods

- Hamiltonian simulation from Trotter-Suzuki decomposition was first proposed by Lloyd (1996).
- In the 2010s, many new ideas appeared. They are interesting, because they come with a better complexity bound than the Trotter-Suzuki method.
- Strangely, many of these algorithms attempt to approximate e^{-iHt} by a *non-unitary operation*.
- We will talk about one such method, known as the LCU(=linear combination of unitaries).

Basic intuition

- What we want: $|\psi\rangle \rightarrow |\psi'\rangle = e^{-iHt} |\psi\rangle$.
- What we do: $|\psi\rangle |0\rangle \rightarrow \sqrt{p} |\widetilde{\psi}'\rangle |Success\rangle + \sqrt{1-p} |Fail\rangle$, where $|\widetilde{\psi}'\rangle \approx |\psi'\rangle$, by applying a linear combination of unitaries (LCU). [Childs and Wiebe (2012)]
- Then we “boost” the success probability to 1.

Measure the 2nd register

┌	Prob p : $ \widetilde{\psi}'\rangle$
	Prob $1-p$: “Fail”

Toy example 1

- Suppose we want to apply an arbitrary 2×2 matrix on a qubit, using a unitary quantum circuit. What can we do?

A : 2×2 matrix

$|\psi\rangle$: 1-qubit state

$$|\psi\rangle \rightarrow A|\psi\rangle / \|A|\psi\rangle\|$$

$$(\| |\psi\rangle \| = \sqrt{\langle \psi | \psi \rangle})$$

Fact: $A = d_z I + d_x X + d_y Y + d_z Z$, $(d_z, d_x, d_y, d_z \in \mathbb{C})$

Register with 2 qubits $|I\rangle, |X\rangle, |Y\rangle, |Z\rangle$
 $|00\rangle, |01\rangle, |10\rangle, |11\rangle$

$$(d_z |I\rangle + d_x |X\rangle + d_y |Y\rangle + d_z |Z\rangle) |\psi\rangle \xrightarrow{U} d_z |I\rangle |\psi\rangle + d_x |X\rangle |X\rangle |\psi\rangle + d_y |Y\rangle |Y\rangle |\psi\rangle + d_z |Z\rangle |Z\rangle |\psi\rangle$$

$$U |I\rangle |\psi\rangle = |I\rangle |\psi\rangle$$

$$U |X\rangle |\psi\rangle = |X\rangle |X\rangle |\psi\rangle$$

$$U |Y\rangle |\psi\rangle = |Y\rangle |Y\rangle |\psi\rangle$$

$$U = I$$

↓

Measure in the basis of $\left\{ \frac{1}{2}(|I\rangle + |X\rangle + |Y\rangle + |Z\rangle), \dots \right\}$

$$U(|z\rangle|\psi\rangle) = |z\rangle z|\psi\rangle$$

$|\phi\rangle$
 \downarrow If I measure $|\phi\rangle$
 $\alpha_z|\psi\rangle + \alpha_x|x\rangle|\psi\rangle + \alpha_y|y\rangle|\psi\rangle + \alpha_z|z\rangle|\psi\rangle$

$$\frac{1}{\sqrt{2}} \left(|z\rangle + |x\rangle + |y\rangle + |z\rangle \right) \left(\alpha_z |I\rangle |\psi\rangle + \alpha_x |X\rangle |\psi\rangle + \alpha_y |Y\rangle |\psi\rangle + \alpha_z |Z\rangle |\psi\rangle \right)$$

Toy example 2

- Suppose we want to apply $e^{-iH\delta t} \approx I - iH\delta t$, where H is a local Hamiltonian. What can we do?

$$H = -t \sum_i (a_i^\dagger a_{i+1} + h.c.) + U \sum_i \hat{n}_i \hat{n}_{i+1}.$$

$$H = -\frac{t}{2} \sum_i (\underbrace{X_i X_{i+1}} + \underbrace{Y_i Y_{i+1}}) + U \sum_i \frac{(Z_i + 1)(Z_{i+1} + 1)}{4}.$$

X_i, Y_i, Z_i

$$H = \sum_P \alpha_P P$$

\leftarrow Tensor product of Pauli operators
 ex) $X_1 \otimes X_2 \otimes Y_3 \otimes I_4 \dots$

$$(|Z\rangle - i\delta t \sum_P \alpha_P |P\rangle) |\Psi\rangle \xrightarrow{U} |Z\rangle |\Psi\rangle - i\delta t \sum_P \alpha_P |P\rangle |\Psi\rangle$$

$$U|p\rangle|\psi\rangle = |p\rangle P|\psi\rangle$$

Hamiltonian Simulation

- Suppose we want to apply $e^{-iH\delta t}$, where H is a local Hamiltonian and the norm of $H\delta t$ is ≈ 0 . What can we do?
- A: Taylor expansion! [Berry, Childs, Cleve, Kohtari, and Somma (2014)]

$$H = -\frac{t}{2} \sum_i (X_i X_{i+1} + Y_i Y_{i+1}) + U \sum_i \frac{(Z_i + 1)(Z_{i+1} + 1)}{4}.$$

$$e^{-iHT} = \left(e^{-\frac{iHT}{n}} \right)^n$$

norm is $O(1)$

$$e^{-\frac{iHT}{n}} = I - \frac{iHT}{n} - \frac{H^2 T^2}{2n^2} + \dots$$

$$= \left(\sum_p \alpha_p P \right) |\psi\rangle \xrightarrow{U} \sum_p \alpha_p |p\rangle P|\psi\rangle$$

$$\begin{array}{l}
 \text{measure } \left\{ \sum_P |P\rangle, \dots \right\} \\
 \text{basis } e^{-\frac{iHt}{\hbar}} |\psi\rangle \\
 \text{measure } \left\{ \sum_P |P\rangle \right\} \downarrow \sum_P d_P P |\psi\rangle \approx
 \end{array}$$

SELECT

- SELECT is an important subroutine used in modern quantum algorithms.
- Think of it as a “lookup table.”
- We will see this subroutine again, later in the course.

$$\{P_1, P_2, \dots, P_n\}$$

$$\begin{array}{l}
 |k\rangle |\psi\rangle \longrightarrow \text{SELECT } |k\rangle |\psi\rangle \\
 \uparrow \\
 \text{Index} \qquad \qquad \qquad = |k\rangle P_k |\psi\rangle
 \end{array}$$

$$U |P\rangle |\psi\rangle = |P\rangle P |\psi\rangle$$

PREPARE

- PREPARE is another important subroutine.
- In the context of the Hamiltonian simulation, this prepares the state that encodes the coefficients of the Hamiltonian.
- We will see this subroutine again, later in the course.

$$\sum_p \alpha_p |p\rangle$$

$$\text{PREPARE } |0 \dots 0\rangle = \sum_{k=1}^n \alpha_k |k\rangle$$

$$\text{SELECT } |k\rangle |\psi\rangle = |k\rangle p_k |\psi\rangle \quad (\{p_1, \dots, p_n\})$$

SELECT + PREPARE

- With these two, we can implement the desired operation.

$$\text{PREPARE } |0 \dots 0\rangle = \sum_{k=1}^n \alpha_k |k\rangle \quad \leftarrow \log_2 n \text{-qubit state}$$

$$\text{SELECT } |k\rangle |\psi\rangle = |k\rangle P_k |\psi\rangle \quad (\{P_1, \dots, P_n\})$$

$$\text{PREPARE}^\dagger \text{ SELECT PREPARE } (|0 \dots 0\rangle |\psi\rangle)$$

$$= \text{PREPARE}^\dagger \text{ SELECT } \left(\sum_{k=1}^n \alpha_k |k\rangle |\psi\rangle \right)$$

$$= \text{PREPARE}^\dagger \sum_{k=1}^n \alpha_k |k\rangle P_k |\psi\rangle$$

$$\langle 0 \dots 0 | \text{PREPARE}^\dagger \sum_{k=1}^n \alpha_k |k\rangle P_k |\psi\rangle = \sum_{k'=1}^n \sum_{k=1}^n \alpha_{k'}^* \alpha_k \langle k' | k \rangle P_k |\psi\rangle$$

$$= \sum_{k=1}^n |\alpha_k|^2 P_k |\psi\rangle$$

$k=1$ —

$$e^{-iHt} \approx \sum_k |k\rangle \langle k| P_k$$

$$H_L = \sum_{i=1}^n z_i z_{i+1} |i\rangle \langle i|^2$$

↓

$$\text{SELECT } |i\rangle |\psi\rangle \\ = |i\rangle z_i z_{i+1} |\psi\rangle$$

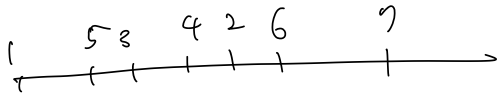
$$\text{PREPARE} = \sum_{i=1}^n |i\rangle \langle i|$$

$$H_L = \sum_{i=1}^n X_i X_{i+1}$$

↓

$$\text{SELECT } |i\rangle |\psi\rangle \\ = |i\rangle X_i X_{i+1} |\psi\rangle$$

$$\text{PREPARE} = \sum_{i=1}^n \frac{1}{\sqrt{n}} |i\rangle$$



\approx

$>$

Why bother with this approach?

- In the Trotter-based Hamiltonian simulation, there is an inevitable $O(\text{poly}(\epsilon^{-1}))$ scaling in the precision ϵ .
- Using the LCU approach, one can get a sub-logarithmic scaling in $1/\epsilon$. [Berry, Childs, Cleve, Kohtari, and Somma (2013, 2014)] $\frac{\log(1/\epsilon)}{\log b_2(1/\epsilon)}$
- This is especially important for quantum chemistry applications, because they tend to require high-precision calculations.

Summary

- In the Trotter-based Hamiltonian simulation, there is an inevitable $O(\text{poly}(\epsilon^{-1}))$ scaling in the precision ϵ .
- Using the LCU approach, one can get a sub-logarithmic scaling in $1/\epsilon$. [Berry, Childs, Cleve, Kohtari, and Somma (2013, 2014)]
- Using SELECT + PREPARE, we can apply the desired unitary with a nonzero probability.
- However, we haven't discussed how to boost this probability. That will come in the next lecture.